

Reconstruction Algorithms for Blind Ptychographic Imaging

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What is Ptychography?

- An unknown **specimen** is illuminated by a **localized illumination function** resulting in an **exit-wave** whose intensity is observed.
- A **ptychography dataset** is a series of these observations, each of which is obtained by shifting the illumination function to a different position relative to the specimen. **Neighbouring illumination regions overlap**.
- Given a ptychographic dataset, the **blind ptychography problem** is to simultaneously reconstruct the relative phase of the specimen and illumination function.

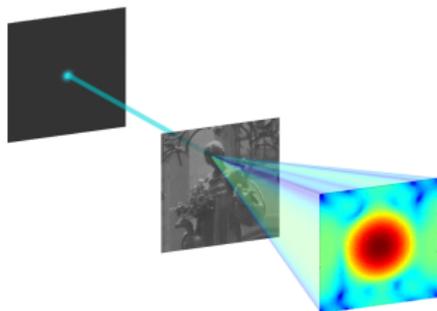


Figure : An illumination function (left), specimen (center), and exit-wave (right).

What is Ptychography?

The mathematical model is:

- $x \in \mathbb{C}^{n \times n}$ is the unknown illumination function,
- $y \in \mathbb{C}^{n \times n}$ is the unknown specimen.
- $\mathbf{z} = (z_1, \dots, z_m) \in (\mathbb{C}^{n \times n})^m$ is an m -tuple of diffraction patterns.
- $S_j : \mathbb{C}^{n \times n} \rightarrow \mathbb{C}^{n \times n}$ is a *shift map* with $S_j(x)$ corresponding to the position of the illumination function for the j^{th} diffraction pattern.
- The elements of the triple (x, y, \mathbf{z}) are related by:

$$S_j(x) \odot y = z_j \quad \text{for } j = 1, 2, \dots, m.$$

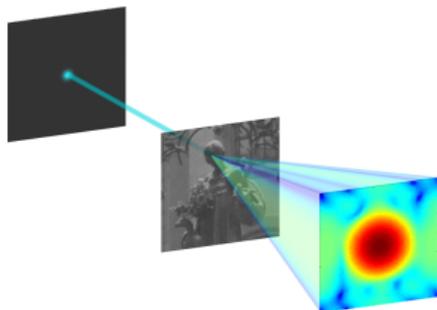


Figure : An example of $S_j(x) \odot y = z_j$ with S_j localising “ x ” to the girl’s head.

What is Ptychography?

In a ptychography experiment we observe the m -matrices

$$b_1, \dots, b_m \in \mathbb{R}_+^{n \times n},$$

where b_j , for $j = 1, 2, \dots, m$, are given by

$$b_j = |\mathcal{F}(z_j)| = |\mathcal{F}(S_j(x) \odot y)|.$$

Here \mathcal{F} is the *2D Fourier transform*, and $|\cdot|$ is taken element-wise.

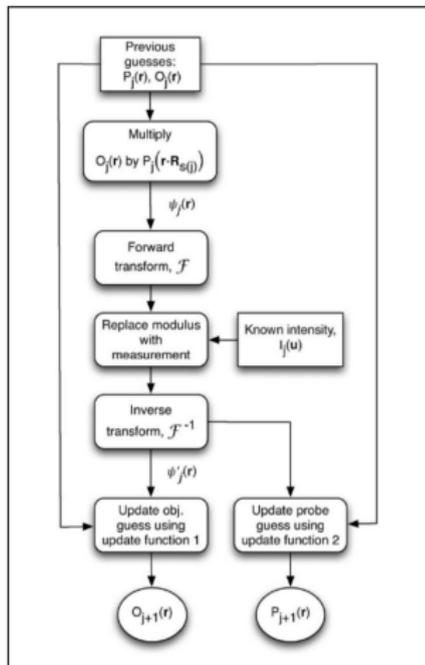
The **blind ptychography problem** is:

Given $b_1, b_2, \dots, b_m \in \mathbb{R}_+^{n \times n}$ reconstruct the triple (x, y, z) .

- Ill-posed, inverse problem with many solutions \Rightarrow hopeless without *a priori* knowledge.

Two Algorithms in the Literature

Maiden & Rodenburg proposed:

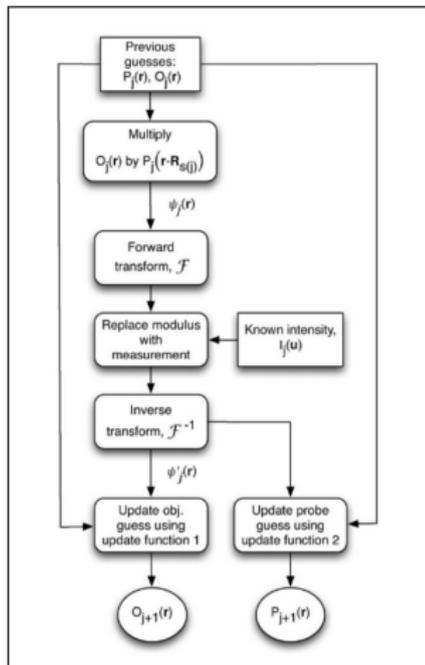


With update functions:

$$O_{j+1}(\mathbf{r}) = O_j(\mathbf{r}) + \alpha \frac{P_j^*(\mathbf{r} - \mathbf{R}_{s(j)})}{|P_j(\mathbf{r} - \mathbf{R}_{s(j)})|_{\max}^2} (\psi'_j(\mathbf{r}) - \psi_j(\mathbf{r})), \quad P_{j+1}(\mathbf{r}) = P_j(\mathbf{r}) + \beta \frac{O_j^*(\mathbf{r} + \mathbf{R}_{s(j)})}{|O_j(\mathbf{r} + \mathbf{R}_{s(j)})|_{\max}^2} (\psi'_j(\mathbf{r}) - \psi_j(\mathbf{r})).$$

Two Algorithms in the Literature

Maiden & Rodenburg proposed:



Thibault *et al.* proposed:

$$\Pi_F(\Psi) : \psi_j \rightarrow \psi_j^F = p_F(\psi_j). \quad (4)$$

$$\Pi_O(\Psi) : \psi_j \rightarrow \psi_j^O(\mathbf{r}) = \hat{P}(\mathbf{r} - \mathbf{r}_j) \hat{O}(\mathbf{r}). \quad (6)$$

$$\Psi_{n+1} = \Psi_n + \Pi_F[2\Pi_O(\Psi_n) - \Psi_n] - \Pi_O(\Psi_n). \quad (9)$$

On computing (6):

$$\hat{O}(\mathbf{r}) = \frac{\sum_j \hat{P}^*(\mathbf{r} - \mathbf{r}_j) \psi_j(\mathbf{r})}{\sum_j |\hat{P}(\mathbf{r} - \mathbf{r}_j)|^2}. \quad (7)$$

$$\hat{P}(\mathbf{r}) = \frac{\sum_j \hat{O}^*(\mathbf{r} + \mathbf{r}_j) \psi_j(\mathbf{r} + \mathbf{r}_j)}{\sum_j |\hat{O}(\mathbf{r} + \mathbf{r}_j)|^2}. \quad (8)$$

In the event the probe \hat{P} is already known, the overlap projection is given by (6), where \hat{O} is computed with Eq. (7). If \hat{P} also needs to be retrieved, both Eqs. (7) and (8) need to be simultaneously solved. **While the system cannot be decoupled analytically, applying the two equations in turns for a few iterations was observed to be an efficient procedure to find the minimum.** Within the reconstruction scheme, initial guesses for \hat{P} and \hat{O} are readily available from the previous iteration—apart

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Our Framework

- Perfectly good algorithmic schemes, which have been shown to work.
- Not clear what (optimisation?) problem the algorithms solve.
 - Cannot be cast as projection-type algorithms solving feasibility problems, although they seem closely related.
- We considered the following **optimisation problem**:

$$\begin{aligned} \min \quad & F(x, y, \mathbf{z}) := \sum_{j=1}^m \|S_j(x) \odot y - z_j\|^2 \\ \text{s.t.} \quad & x \in X = \{x : \|x\|_\infty \leq M_x, x|_{\mathbb{I}_x} = 0\}, \\ & y \in Y = \{y : \|y\|_\infty \leq M_y\}, \\ & \mathbf{z} \in Z = \{z : |\mathcal{F}(z_j)| = b_j \text{ for } j = 1, 2, \dots, m\}, \end{aligned} \quad (\text{P})$$

where $M_x, M_y \in \mathbb{R}$ are bounds, and \mathbb{I}_x is an index set (**support** of x).

- Separable constraint sets coupled through a **“nice”** objective function.
- (P) is equivalent to the formally unconstrained problem:

$$\min \Psi(x, y, \mathbf{z}) := F(x, y, \mathbf{z}) + \iota_X(x) + \iota_Y(y) + \iota_Z(\mathbf{z}),$$

where $\iota_C(w)$ is the **indicator function** of the set C which takes the value 0 if $w \in C$, and $+\infty$ if $w \notin C$.

A Naive Approach: Alternating Minimisation

Alternating Minimisation Algorithm (over three blocks):

Initialization. Choose $(x^0, y^0, z^0) \in X \times Y \times Z$.

General Step. ($k = 0, 1, \dots$)

1. Select $x^{k+1} \in \arg \min_{x \in X} F(x, y^k, z^k),$

2. Select $y^{k+1} \in \arg \min_{y \in Y} F(x^{k+1}, y, z^k),$

3. Select $z^{k+1} \in \arg \min_{z \in Z} F(x^{k+1}, y^{k+1}, z).$

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3. Select $z^{k+1} \in \arg \min_{z \in Z} F(x^{k+1}, y^{k+1}, z).$

What's involved? Roughly speaking, to compute Step 1 we must minimise the terms of the form $\|S_j(x) \odot y^k - z_j^k\|^2$. If this is zero, then

$$S_j(x) \odot y^k = z_j^k \implies S_j(x) = z_j^k \oslash y_k \implies x = S_j^{-1}(z_j^k \oslash y_k).$$

- Inverting S_j is stable (just “un-shift”).
- Division by y_k is unstable (divide by zero)
- Similar observations apply to Step 2.
- Step 3 is unstable but there are regularisation schemes in the literature.

PHeBIE: Proximal Block Implicit-Explicit Algorithm

We **regularise** doing the following to F :

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- Linearising Steps 1 & 2. For Step 1, we have

$$F(x, y^k, \mathbf{z}^k) \longrightarrow F(x^k, y^k, \mathbf{z}^k) + \langle x - x^k, \nabla_x F(x^k, y^k, \mathbf{z}^k) \rangle.$$

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- Adding a **proximal term** to Steps 1, 2 & 3. For Step 1, such a term looks like

$$+ \frac{\alpha^k}{2} \|x - x^k\|^2,$$

where $\alpha^k > 0$ is dependent on y^k and \mathbf{z}^k .

PHeBE: Proximal Block Implicit-Explicit Algorithm

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The regularised version of Step 1 becomes:

$$x^{k+1} \in \arg \min_{x \in X} \left\{ \langle x - x^k, \nabla_x F(x^k, y^k, \mathbf{z}^k) \rangle + \frac{\alpha^k}{2} \|x - x^k\|^2 \right\}.$$

The other steps are similar.

PHeBE: Proximal Block Implicit-Explicit Algorithm

Proximal Block Implicit-Explicit Algorithm:

Initialization. Choose $\alpha, \beta > 1$, $\gamma, \eta_x, \eta_y > 0$, $(x^0, y^0, z^0) \in X \times Y \times Z$.

General Step. ($k = 0, 1, \dots$)

1. Set $\alpha^k = \alpha \max\{L_x(y^k, z^k), \eta_x\}$ and select

$$x^{k+1} \in P_X \left(x^k - \frac{2}{\alpha^k} \sum_{j=1}^m S_j^{-1}(\bar{y}^k) \odot S_j^{-1}(y^k - z_j^k) \right),$$

2. Set $\beta^k = \beta \max\{L_y(x^{k+1}, z^k), \eta_y\}$ and select

$$y^{k+1} \in P_Y \left(y^k - \frac{2}{\beta^k} \sum_{j=1}^m S_j(\bar{x}^{k+1}) \odot (S_j(x^{k+1}) - z_j^k) \right),$$

3. Select, for $j = 1, 2, \dots, m$,

$$z_j^{k+1} \in P_Z \left(\left[\frac{2}{2 + \gamma_k} S_j(x^{k+1}) \odot y^{k+1} + \frac{\gamma_k}{2 + \gamma_k} z_j^k \right]_{j=1}^m \right).$$

Here $L_x(y^k, z^k)$ (resp $L_y(x^{k+1}, z^k)$) denotes the partial Lipschitz constant of $\nabla_x F(x, y^k, z^k)$ (resp. $\nabla_y F(x^{k+1}, y, z^k)$), and the projection onto a set C is given by

$$P_C(w) := \arg \min_{u \in C} \|u - w\|^2.$$

PHeBIE: Convergence Theorem

Theorem (Hesse–Luke–Sabach–T, 2015)

Let $\{(x^k, y^k, z^k)\}_{k \in \mathbb{N}}$ be a sequence generated by PHeBIE for blind ptychography problem. Then the following hold.

- ① The sequence $\{(x^k, y^k, z^k)\}_{k \in \mathbb{N}}$ has finite length. That is,

$$\sum_{k=1}^{\infty} \|(x^{k+1}, y^{k+1}, z^{k+1}) - (x^k, y^k, z^k)\|^2 < \infty.$$

- ② The sequence $\{(x^k, y^k, z^k)\}_{k \in \mathbb{N}}$ converges to point (x^*, y^*, z^*) which is a critical point of the function Ψ . That is,

$$0 \in \partial\Psi(x, y, z) = \nabla F(x^*, y^*, z^*) + \partial_{\ell_X}(x^*) + \partial_{\ell_Y}(y^*) + \partial_{\ell_Z}(z^*),$$

where $\partial(\cdot)$ denotes the **limiting Fréchet subdifferential**.

For u in the domain of f , the **limiting Fréchet subdifferential** is given by

$$\partial f(u) := \left\{ v : \exists u^k \rightarrow u, f(u^k) \rightarrow f(u), v^k \rightarrow v, v^k \in \widehat{\partial} f(u^k) \right\}, \text{ where } \widehat{\partial} f(u) = \left\{ v : \liminf_{\substack{w \neq u \\ w \rightarrow u}} \frac{f(w) - f(u) - \langle v, w - u \rangle}{\|w - u\|} \geq 0 \right\}.$$

PHeBIE: Example

PHeBIE: Convergence Theorem (cont.)

Proof Sketch.

The proof has three steps:

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- 1 (Sufficient decrease) Use properties of the algorithm to establish that the sequence $\{F(x^k, y^k, z^k)\}_{k \in \mathbb{N}}$ is decreasing, converges to some $F^* > -\infty$, and that

$$\sum_{k=1}^{\infty} \|(x^{k+1}, y^{k+1}, z^{k+1}) - (x^k, y^k, z^k)\|^2 < \infty.$$

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- 2 (Subdifferential bound) Use properties of the algorithm to show that

$$\|w^{k+1}\| \leq \kappa \|(x^{k+1}, y^{k+1}, \mathbf{z}^{k+1}) - (x^k, y^k, \mathbf{z}^k)\|,$$

for some $w^{k+1} \in \partial\Psi(x^{k+1}, y^{k+1}, \mathbf{z}^{k+1})$ and $\kappa > 0$.

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- 3 To establish convergence of $\{(x^k, y^k, \mathbf{z}^k)\}_{k \in \mathbb{N}}$ to a critical point, we appeal to a result of *Bolte, Sabach & Teboulle, 2014*. Here the important ingredient is that Ψ satisfied the **Kurdyka–Łojasiewicz (KL) Property** which gives Cauchyness of $\{(x^k, y^k, \mathbf{z}^k)\}_{k \in \mathbb{N}}$.



The Kurdyka-Łojasiewicz (KL) Property

Let $f : \mathbb{R}^d \rightarrow (-\infty, +\infty]$ be proper and lsc. For $\eta \in (0, +\infty]$ define

$$\mathcal{C}_\eta \equiv \{\phi : [0, \eta) \rightarrow \mathbb{R}_+ : \phi(0) = 0, \phi'(s) > 0 \text{ for all } s \in (0, \eta)\}.$$

The function f is said to have the **KL property** at $\bar{u} \in \text{dom } \partial f$ if there exists a neighbourhood U of \bar{u} and a function $\varphi \in \mathcal{C}_\eta$, such that, for all

$$u \in \{u \in U : f(\bar{u}) < f(u) < f(\bar{u}) + \eta\},$$

it holds that

$$\varphi'(f(u) - f(\bar{u})) \text{dist}(0, \partial f(u)) \geq 1.$$

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Why is the KL property useful? It holds for many important nonconvex optimisation problems.

In particular, any proper, lsc, **semi-algebraic** function satisfies the KL property everywhere in its domain.

Exploiting Block Structure

The present algorithm:

- Alternatively minimizes w.r.t. three blocks: X , Y and Z .
- At each iteration the “step-size” is inversely proportional to a **parital Lipschitz constant**. For instance,

$$L_x(y^k, z^k) = 2 \left\| \sum_{j=1}^m S_j^*(\bar{y}^k \odot y^k) \right\|_{\infty}.$$

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When X or Y has separable structure, they can be decomposed as

$$X \equiv X_1 \times X_2 \times \cdots \times X_N, \quad Y \equiv Y_1 \times Y_2 \times \cdots \times Y_M.$$

- Algorithm variant: alternatively minimizes w.r.t. $(N + M + 1)$ -blocks.
- The j th “sub-block”, X_j , has partial Lipschitz constant

$$2 \left\| \left(\sum_{j=1}^m S_j^*(\bar{y}^k \odot y^k) \right) \right\|_{X_j} \Big|_{\infty}.$$

- ★ More sub-blocks \rightarrow small sub-blocks \rightarrow small constant \rightarrow **larger step-size**.
- Sub-blocks can be updated **sequentially** or **in parallel**. In both cases, an analogous convergence theorem holds (see the paper for details).

Explaining Maiden & Rodenburg

In our framework, we can interpret Maiden & Rodenburg's algorithm as alternating minimisation w.r.t. the blocks (in order)

$$X_1, Y_1, Z, X_2, Y_2, Z, X_3, Y_3, Z, \dots, X_1, Y_1, Z, \dots$$

where sub-blocks X_j and Y_j correspond to restrictions related to the support of the j th exit wave.

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where sub-blocks X_j and Y_j correspond to restrictions related to the support of the j th exit wave.

Moreover, Maiden & Rodenburg's update rules:

$$O_{j+1}(\mathbf{r}) = O_j(\mathbf{r}) + \alpha \frac{P_j^*(\mathbf{r} - \mathbf{R}_{s(j)})}{|P_j(\mathbf{r} - \mathbf{R}_{s(j)})|_{\max}^2} (\psi'_j(\mathbf{r}) - \psi_j(\mathbf{r})), \quad P_{j+1}(\mathbf{r}) = P_j(\mathbf{r}) + \beta \frac{O_j^*(\mathbf{r} + \mathbf{R}_{s(j)})}{|O_j(\mathbf{r} + \mathbf{R}_{s(j)})|_{\max}^2} (\psi'_j(\mathbf{r}) - \psi_j(\mathbf{r})).$$

are what is obtain by by taking $X = Y = \mathbb{C}^{m \times m}$. Note that the "normalisations" are precisely the **partial Lipschitz constants!**

Explaining Thibault *et al.*

If it were not a **difference-map** algorithm, Thibault *et al.* could be interpreted as the two block alternating proximal linearisation minimisation of:

Initialization. Choose $(x^0, y^0, z^0) \in X \times Y \times Z$.

General Step. ($k = 0, 1, \dots$)

1. Select $(x^{k+1}, y^{k+1}) \in \arg \min_{(x,y) \in X \times Y} F(x, y, z^k),$

2. Select $z^{k+1} \in \arg \min_{z \in Z} F(x^{k+1}, y^{k+1}, z).$

However Step. 1 is not so easy to solve:

$$\hat{O}(\mathbf{r}) = \frac{\sum_j \hat{P}^*(\mathbf{r} - \mathbf{r}_j) \psi_j(\mathbf{r})}{\sum_j |\hat{P}(\mathbf{r} - \mathbf{r}_j)|^2}, \quad (7)$$

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2. Select $z^{k+1} \in \arg \min_{z \in Z} F(x^{k+1}, y^{k+1}, z).$

So the actual computation is the following **heuristic approximation** to 1.:

Input. $(\hat{x}^0, \hat{y}^0) := (x^k, y^k) \in X \times Y, z^k$, and $L \in \mathbb{N}$.

General Step. ($l = 0, 1, \dots, (L-1)$)

1a. Select $\hat{x}^{l+1} \in \arg \min_{x \in X} F(x, \hat{y}^l, z^k),$

1b. Select $\hat{y}^{l+1} \in \arg \min_{y \in Y} F(\hat{x}^{l+1}, y, z^k),$

In our framework, this is alternating minimisation w.r.t. the blocks (in order)

$X, Y, X, Y, \dots, X, Y, Z, X, Y, X, Y, \dots, X, Y, Z, \dots$

Concluding Remarks and Ongoing Work

In summary:

- We presented a ptychography algorithm with a clear mathematical framework.
- Under practically verifiable assumption, it is provably convergent to critical points of a function Ψ .
- The flexibility of the framework allows interpretation of current state-of-the-art ptychography algorithms.

Ongoing and future work:

- Is there a useful characterisation of the critical points of Ψ ?
- Can the algorithm structure be exploited on specific architectures?

For further details see:

Proximal Heterogeneous Block Implicit-Explicit Method and Application to Blind Ptychographic Diffraction Imaging with R. Hesse, D.R. Luke and S. Sabach. *SIAM J. on Imaging Sciences*, 8(1):426–457 (2015).